# A New Look at Hamiltonian Walks

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#### ABSTRACT

Let G be a connected graph of order n. A Hamiltonian walk of G is a closed spanning walk of minimum length in G. For a cyclic ordering  $s: v_1, v_2, \dots, v_n, v_{n+1} = v_1$  of V(G), let d(s) = $\sum_{i=1}^{n} d(v_i, v_{i+1})$ , where  $d(v_i, v_{i+1})$  is the distance between  $v_i$ and  $v_{i+1}$  for  $1 \leq i \leq n$ . Then the Hamiltonian number h(G) of G is defined as  $h(G) = \min \{d(s)\}$ , where the minimum is taken over all cyclic orderings s of V(G). It is shown that h(G) is the length of a Hamiltonian walk in G. Thus h(G) = n if and only if G is a Hamiltonian graph. It is also shown that h(G) = 2n - 2 if and only if G is a tree. Moreover, for every pair n, k of integers with  $3 \le n \le k \le 2n-2$ , there exists a connected graph G of order n having h(G) = k. The upper Hamiltonian number is defined as  $h^+(G) = \max \{ d(s) \}$ , where the maximum is taken over all cyclic orderings s of V(G). We show, for a connected graph G of order  $n \geq 3$ , that  $h(G) = h^+(G)$  if and only if  $G = K_n$  or  $G = K_{1,n-1}$ . We also study the upper Hamiltonian number of a tree and present bounds for the upper Hamiltonian number of a connected graph in terms of its order.

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### 1 Introduction

In [6] Goodman and Hedetniemi introduced the concept of a Hamiltonian walk in a connected graph G, defined as a closed spanning walk of minimum length in G. They denoted the length of a Hamiltonian walk in G by h(G). Therefore, for a connected graph G of order  $n \ge 3$ , it follows that h(G) = n if and only if G is Hamiltonian. Among the results obtained by Goodman and Hedetniemi are the following.

**Theorem A** If T is a tree of order n, then h(T) = 2n - 2.

It is immediate that  $h(G) \leq h(H)$  for each connected spanning subgraph H of a (connected) graph G. As a consequence of Theorem A, we can state the following.

**Theorem B** For every connected graph G of order n,

 $n \le h(G) \le 2n - 2.$ 

**Theorem C** If G is a k-connected graph of order n having diameter d, then

$$h(G) \le 2n - \left\lfloor \frac{k}{2} \right\rfloor (2d-2) - 2.$$

**Theorem D** Let G be a connected graph having blocks  $B_1, B_2, \ldots, B_k$ . Then the union of the edges in a Hamiltonian walk for each of the blocks  $B_i$ forms a Hamiltonian walk for G and, conversely, the edges in a Hamiltonian walk of G that belong to  $B_i$  form a Hamiltonian walk in  $B_i$ .

Theorem D implies that the topic of Hamiltonian walks can be restricted to 2-connected graphs. Hamiltonian walks were studied further in [1, 2, 3, 5, 8, 9]. A well-known sufficient condition for a graph G to be Hamiltonian is due to Ore [7].

**Theorem E** A graph G of order  $n \ge 3$  is Hamiltonian if  $\deg u + \deg v \ge n$ for every pair u, v of nonadjacent vertices of G.

This theorem can be restated in terms of the parameter h(G).

**Theorem F** Let G be a graph of order  $n \ge 3$ . Then h(G) = n if  $\deg u + \deg v \ge n$  for every pair u, v of nonadjacent vertices of G.

Bermond [3] obtained the following generalization of Theorem F.

**Theorem G** Let G be a connected graph G of order  $n \ge 3$  and let k be an integer with  $0 \le k \le n-2$ . If  $\deg u + \deg v \ge n-k$  every pair u, v of nonadjacent vertices of G, then  $h(G) \le n+k$ . In this paper, we refer to the book [4] for graph theory notation and terminology not described here.

### 2 The Hamiltonian Number of a Graph

Of course, a Hamiltonian graph G contains a spanning cycle  $C: v_1, v_2, \dots, v_n, v_{n+1} = v_1$ , where then  $v_i v_{i+1} \in E(G)$  for  $1 \leq i \leq n$ . Thus Hamiltonian graphs of order  $n \geq 3$  are those graphs for which there is a cyclic ordering  $v_1, v_2, \dots, v_n, v_{n+1} = v_1$  of V(G) such that  $\sum_{i=1}^n d(v_i, v_{i+1}) = n$ , where  $d(v_i, v_{i+1})$  is the distance between  $v_i$  and  $v_{i+1}$  for  $1 \leq i \leq n$ . For a connected graph G of order  $n \geq 3$  and a cyclic ordering  $s: v_1, v_2, \dots, v_n, v_{n+1} = v_1$  of V(G), we define the number d(s) by

$$d(s) = \sum_{i=1}^{n} d(v_i, v_{i+1}).$$

Therefore,  $d(s) \ge n$  for each cyclic ordering s of V(G). The Hamiltonian number  $h^*(G)$  of G is defined by

$$h^*(G) = \min\left\{d(s)\right\},\,$$

where the minimum is taken over all cyclic orderings s of V(G). Consider the graph  $G = K_{2,3}$  of Figure 1. For the cyclic orderings

 $s_1: v_1, v_2, v_3, v_4, v_5, v_1$  and  $s_2: v_1, v_3, v_2, v_4, v_5, v_1$ 

of V(G), we see that  $d(s_1) = 8$  and  $d(s_2) = 6$ . Since G is a non-Hamiltonian graph of order 5 and  $d(s_2) = 6$ , it follows that  $h^*(G) = 6$ .

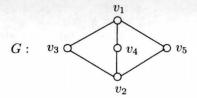


Figure 1: A graph G with  $h^*(G) = 6$ 

We now see that there is an alternative way to define the length h(G) of a Hamiltonian walk in G. Denote the length of a walk W by L(W).

**Proposition 2.1** For every connected graph G,

$$h^*(G) = h(G).$$

**Proof.** First, we show that  $h(G) \leq h^*(G)$ . Let  $s: v_1, v_2, \dots, v_n, v_{n+1} = v_1$  be a cyclic ordering of V(G) for which  $d(s) = h^*(G)$ . For each integer i with  $1 \leq i \leq n$ , let  $P_i$  be a  $v_i - v_{i+1}$  geodesic in G. Thus  $L(P_i) = d(v_i, v_{i+1})$ . The union of the paths  $P_i$  forms a closed spanning walk W in G. Therefore,

$$h(G) \le L(W) = \sum_{i=1}^{n} L(P_i) = \sum_{i=1}^{n} d(v_i, v_{i+1}) = d(s) = h^*(G).$$

Next, we show that  $h^*(G) \leq h(G)$ . Let W be a Hamiltonian walk in G. Therefore, L(W) = h(G). Suppose that  $W: x_1, x_2, \ldots, x_N, x_1$ , where then  $N \geq n$ . Define  $v_1 = x_1$  and  $v_2 = x_2$ . For  $3 \leq i \leq n$ , define  $v_i$  to be  $x_{j_i}$ , where  $j_i$  is the smallest positive integer such that  $x_{j_i} \notin \{v_1, v_2, \ldots, v_{i-1}\}$ . Then  $s: v_1, v_2, \ldots, v_n, v_{n+1} = v_1$  is a cyclic ordering of V(G). For each i with  $1 \leq i \leq n$ , let  $W_i$  be the  $v_i - v_{i+1}$  subwalk of W and so  $d(v_i, v_{i+1}) \leq L(W_i)$ . Since

$$h^*(G) \le \sum_{i=1}^n d(v_i, v_{i+1}) \le \sum_{i=1}^n L(W_i) = L(W) = h(G),$$

we have the desired result.

As a consequence of Proposition 2.1, we henceforth denote the Hamiltonian number of a graph G by h(G), which is then the length of a Hamiltonian walk in G.

By Theorem A, if T is a tree of order n, then h(T) = 2n - 2. We now show that the converse of this statement holds as well. To do this, we first state a lemma.

**Lemma 2.2** If G is a connected graph such that  $\delta(G) \ge 2$  and  $\Delta(G) \ge 3$ , then G contains two distinct cycles C and C' such that  $V(C) \ne V(C')$ .

**Theorem 2.3** Let G be a connected graph of order n. Then h(G) = 2n - 2 if and only if G is a tree.

**Proof.** By Theorem A, it suffices to show that if G is a connected graph of order  $n \ge 3$  that is not a tree, then h(G) < 2n - 2. We proceed by induction on n. Since  $h(K_3) = 3$ , the result holds for n = 3. Suppose that h(F) < 2(n-1) - 2 = 2n - 4 for all connected graphs F of order  $n - 1 \ge 3$  that are not trees. Let G be a connected graph of order  $n \ge 4$  that is not a tree. Since  $h(C_n) = n < 2n - 2$ , we may assume that  $G \neq C_n$ .

We claim that G contains a vertex u such that G - u is a connected subgraph of G that is not a tree. If G contains cut-vertices, then there is a vertex u that is a non-cut-vertex of an end-block that has the desired property. So we may assume that G is 2-connected and so  $\delta(G) \geq 2$ . By Lemma 2.2, G contains two distinct cycles C and C' with  $V(C) \neq V(C')$ . Thus if u is a vertex that belongs to one of C and C' but not the other, then G - u is a connected subgraph of G that is not a tree. By the induction hypothesis, h(G - u) < 2(n - 1) - 2 = 2n - 4. Let

$$s_0: v_1, v_2, \ldots, v_{n-1}, v_1$$

be a cyclic ordering of V(G-u) with  $d(s_0) = h(G-u) < 2n-4$ . Suppose that u is adjacent to the vertex  $v_i$ , where  $1 \le i \le n-1$ . Define the cyclic ordering  $s'_0$  of V(G) from  $s_0$  by

$$s'_0: v_1, v_2, \ldots, v_i, u, v_{i+1}, \ldots, v_{n-1}, v_1.$$

Since  $d(v_i, u) = 1$ , it follows by the triangle inequality that

$$d(u, v_{i+1}) \leq 1 + d(v_i, v_{i+1}).$$

Therefore,

$$d(s'_0) = d(s_0) - d(v_i, v_{i+1}) + d(v_i, u) + d(u, v_{i+1})$$
  

$$\leq d(s_0) - d(v_i, v_{i+1}) + 1 + [1 + d(v_i, v_{i+1})]$$
  

$$= d(s_0) + 2 < (2n - 4) + 2 = 2n - 2.$$

Therefore,  $h(G) \leq d(s'_0) < 2n - 2$ , as desired.

By Theorem B, if G is a connected graph G of order n, then  $n \le h(G) \le 2n-2$ . Next we show that every pair k, n of integers with  $3 \le n \le k \le 2n-2$  is realizable as the Hamiltonian number and the order of some connected graph. In order to do this, we first present a known result, which is a consequence of Theorem D (see [6]).

**Corollary H** Let G be a connected graph having blocks  $B_1, B_2, \ldots, B_k$ . Then

$$h(G) = \sum_{i=1}^{k} h(B_i).$$

In particular, every bridge of G appears twice in every Hamiltonian walk of G.

**Proposition 2.4** For every pair n, k of integers with  $3 \le n \le k \le 2n-2$ , there exists a connected graph G of order n having h(G) = k.

**Proof.** For k = n, let G be a Hamiltonian graph of order n; while for k = 2n - 2, let G be a tree of order n. For n < k < 2n - 2, let  $k = n + \ell$ , where  $1 \leq \ell \leq n - 3$ . Now let G be the graph obtained from a cycle  $C_{n-\ell}: u_1, u_2, \ldots, u_{n-\ell}, u_1$  and a path  $P_\ell: v_1, v_2, \ldots, v_\ell$  by joining  $u_1$  to  $v_1$ . Since  $C_{n-\ell}$  is a block of G and any edge not on  $C_{n-\ell}$  is a bridge of G, it then follows by Corollary H that

$$h(G) = h(C_{n-\ell}) + 2\ell = (n-\ell) + 2\ell = n + \ell = k,$$

as desired.

## 3 The Upper Hamiltonian Number of a Graph

We saw for the graph G of Figure 1 that there are cyclic orderings  $s_1$  and  $s_2$  of V(G) such that  $d(s_1) = 8$  and  $d(s_2) = 6$ . Indeed, it is not difficult to see that for *every* cyclic ordering s of V(G), either d(s) = 6 or d(s) = 8.

For a connected graph G, we define the upper Hamiltonian number  $h^+(G)$  by

$$h^+(G) = \max\left\{d(s)\right\},\,$$

where the maximum is taken over all cyclic orderings s of V(G). From our remarks above, it follows that  $h^+(K_{2,3}) = 8$ . As an illustration, we now establish the upper Hamiltonian numbers of the hypercubes.

**Proposition 3.1** For each integer  $n \ge 2$ ,

$$h^+(Q_n) = 2^{n-1}(2n-1).$$

**Proof.** First, we show that  $h^+(Q_n) \leq 2^{n-1}(2n-1)$ . Let s be an arbitrary cyclic ordering of  $V(Q_n)$  with  $d(s) = h^+(Q_n)$ . Since diam  $Q_n = n$  and for each vertex v in  $Q_n$ , there is exactly one vertex in  $Q_n$  whose distance from v is n, it follows that there are at most  $2^{n-1}$  terms in d(s) equal to n. Consequently, each of the remaining  $2^{n-1}$  terms in d(s) is at most n-1. Thus

$$d(s) \le 2^{n-1}n + 2^{n-1}(n-1) = 2^{n-1}(2n-1),$$

and so  $h^+(Q_n) \leq 2^{n-1}(2n-1)$ .

Next we show that  $h^+(Q_n) \ge 2^{n-1}(2n-1)$ . Since the result is true for  $Q_2$ , we may assume that  $n \ge 3$ . Let  $G = Q_n$ . Then G consists of two disjoint copies  $G_1$  and  $G_2$  of  $Q_{n-1}$ , where corresponding vertices of  $G_1$  and  $G_2$  are adjacent. For each vertex v of G, there is a unique vertex  $\overline{v}$  of G such that  $d(v,\overline{v}) = n = \operatorname{diam} Q_n$ . Necessarily, exactly one of v and  $\overline{v}$  belongs to  $G_1$  for each vertex v of G. It is well-known that  $Q_n$  is Hamiltonian for  $n \ge 2$ . Let  $C: v_1, v_2, \ldots, v_{2^{n-1}}, v_{2^{n-1}+1} = v_1$  be a Hamiltonian cycle in  $G_1$ . Now define a cyclic ordering s of V(G) by

$$s: v_1, \overline{v}_1, v_2, \overline{v}_2, \ldots, v_{2^{n-1}}, \overline{v}_{2^{n-1}}, v_1.$$

Since  $d(v_i, \overline{v}_i) = n$  and  $d(v_i, v_{i+1}) = 1$  for  $1 \le i \le 2^{n-1}$ , it follows by the triangle inequality that

$$n = d(v_i, \overline{v}_i) \le d(v_i, v_{i+1}) + d(v_{i+1}, \overline{v}_i) = 1 + d(v_{i+1}, \overline{v}_i).$$

Thus  $d(v_{i+1}, \overline{v}_i) \ge n-1$ , which implies that  $d(v_{i+1}, \overline{v}_i) = n-1$ . Hence

$$h^+(Q_n) \ge d(s) = 2^{n-1}n + 2^{n-1}(n-1) = 2^{n-1}(2n-1),$$

as desired.

Obviously,  $h(G) \leq h^+(G)$  for every connected graph G. For each integer  $n \geq 3$ , there are only two graphs G of order n for which  $h(G) = h^+(G)$ .

**Theorem 3.2** Let G be a connected graph of order  $n \ge 3$ . Then

 $h(G) = h^+(G)$  if and only if  $G = K_n$  or  $G = K_{1,n-1}$ .

**Proof.** If  $G = K_n$ , then certainly d(s) = n for every cyclic ordering s of V(G); while if  $G = K_{1,n-1}$ , then d(s) = 2n - 2 for every cyclic ordering s of V(G). Thus  $h(G) = h^+(G)$  if  $G = K_n$  or  $G = K_{1,n-1}$ .

For the converse, suppose that G is a connected graph of order  $n \geq 3$  such that  $G \neq K_n, K_{1,n-1}$ . We show that  $h(G) \neq h^+(G)$ . Let diam G = d. Since  $G \neq K_n$ , it follows that  $d \geq 2$ . We consider two cases, according to whether  $d \geq 3$  or d = 2.

Case 1.  $d \ge 3$ . Let  $v_1$  and  $v_{d+1}$  be vertices of G such that  $d(v_1, v_{d+1}) = d$ and let  $P: v_1, v_2, \ldots, v_{d+1}$  be a  $v_1 - v_{d+1}$  geodesic in G. Let W = V(G) - V(P). If  $W \neq \emptyset$ , then let  $W = \{w_1, w_2, \ldots, w_\ell\}$ , where  $\ell = n - d - 1$ . Define a cyclic ordering s of V(G) by

$$s: v_1, v_2, v_3, \dots, v_{d+1}, v_1 \text{ or}$$
 (1)

$$s: v_1, v_2, v_3, \dots, v_{d+1}, w_1, w_2, \dots, w_{\ell}, v_1,$$
 (2)

according to whether  $W = \emptyset$  or  $W \neq \emptyset$ . Let s' be the cyclic ordering of V(G) obtained from s by interchanging the locations of  $v_2$  and  $v_3$  in s; that is,

$$s': v_1, v_3, v_2, v_4, \dots, v_{d+1}, v_1$$
 (3)

or 
$$s': v_1, v_3, v_2, v_4, \dots, v_{d+1}, w_1, w_2, \dots, w_\ell, v_1,$$
 (4)

according to whether  $W = \emptyset$  or  $W \neq \emptyset$ . In either case, d(s') = d(s) + 2 and so  $h(G) \neq h^+(G)$ .

Case 2. d = 2. Since G is not a star, it follows that G is not a tree. Thus the girth  $g(G) = k \ge 3$ . Assume first that k = 3. Since G is connected and  $G \ne K_n$ , there exists a set U of four vertices of G such that  $\langle U \rangle = K_4 - e$  or  $\langle U \rangle$  is a triangle with a pendant edge. Therefore, we may assume, without loss of generality, that G contains one of the graphs  $F_1$  and  $F_2$  in Figure 2 as an induced subgraph. In either case, define the cyclic orderings s and s' as

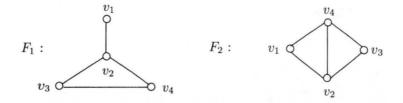


Figure 2: Induced subgraphs  $F_1$  and  $F_2$  of G

described in (1) (or (2)) and (3) (or (4)), respectively. Then d(s') = d(s) + 1and so  $h(G) \neq h^+(G)$ .

If  $k \ge 4$ , then let  $C: v_1, v_2, \ldots, v_k, v_1$  be an induced cycle of G and let  $V(G) - V(C) = \{w_1, w_2, \ldots, w_\ell\}$  if  $\ell = n - k > 0$ . Define the cyclic orderings s and s' of V(G) as in (1) (or (2)) and (3) (or (4)), respectively. Since d(s') = d(s) + 2, it follows that  $h(G) \ne h^+(G)$ .

# 4 Bounds for the Upper Hamiltonian Number of a Graph

First, we observe that if  $s: v_1, v_2, \ldots, v_n, v_{n+1} = v_1$  is any cyclic ordering of the vertex set of a connected graph, then for each vertex  $v_i$   $(1 \le i \le n)$ , both  $d(v_{i-1}, v_i) \le e(v_i)$  and  $d(v_i, v_{i+1}) \le e(v_i)$ , where the subscripts are expressed as integers modulo n and  $e(v_i)$  is the eccentricity of  $v_i$  (the distance from  $v_i$  to a vertex farthest from  $v_i$ ). Thus, If G is a connected graph of order  $n \ge 3$  and  $V(G) = \{v_1, v_2, \ldots, v_n\}$ , then

$$h^+(G) \le \sum_{i=1}^n e(v_i).$$

Since the eccentricity of a vertex in G is at most the diameter diam G of G (the largest distance between two vertices of G), we have the following.

**Proposition 4.1** If G is a connected graph of order  $n \ge 3$  and diameter d, then

$$h^+(G) \le nd.$$

The upper bound in Proposition 4.1 is sharp. For example, consider the odd cycle  $C_{2k+1}: v_1, v_2, \ldots, v_{2k+1}, v_1$ , where  $k \ge 1$ . Since diam  $C_{2k+1} = k$ , it follows by Proposition 4.1 that  $h^+(C_{2k+1}) \le k(2k+1)$ . On the other hand, let

 $s: v_1, v_{k+1}, v_{2k+1}, v_{3k+1}, \dots, v_{(2k)k+1}, v_{(2k+1)k+1} = v_1,$ 

where each subscript is expressed modulo 2k + 1 as one of the integers  $1, 2, \ldots, 2k + 1$ . Since k and 2k + 1 are relatively prime, s is a cyclic ordering of  $V(C_{2k+1})$ . Since

$$d(s) = \sum_{i=1}^{n} d(v_i, v_{i+1}) = k(2k+1),$$

we have the following result.

**Proposition 4.2** For every integer  $k \ge 1$ , let n = 2k+1. Then  $h^+(C_n) = nd$ , where  $d = \operatorname{diam} C_{2k+1}$ .

Therefore, the upper bound in Proposition 4.1 is attained for odd cycles. The situation for even cycles is far less clear. For every integer  $k \ge 2$ , we know that  $h^+(C_{2k}) \ge 2k^2 - 2k + 2$ . Indeed, we state the following.

#### **Conjecture 4.3** For every integer $k \ge 2$ , $h^+(C_{2k}) = 2k^2 - 2k + 2$ .

Next, we study the upper Hamiltonian number of a tree. For each edge e of a tree T, we define the *component number* cn(e) of e as the minimum order of a component of T - e. For example, the edge  $e_3$  of the tree T of Figure 3(a) has component number 3 since the order of the smaller component of  $T - e_3$  is 3. Each edge of this tree is labeled with its component number in Figure 3(b).

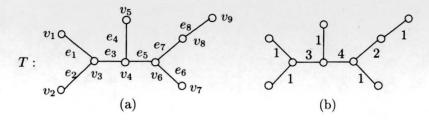


Figure 3: Component numbers of edges

We now present an upper bound for the upper Hamiltonian number of a tree.

**Lemma 4.4** Let T be a tree of order n with  $E(T) = \{e_1, e_2, \dots, e_{n-1}\}$ . Then

$$h^+(T) \le 2\sum_{i=1}^{n-1} \operatorname{cn}(e_i).$$

**Proof.** Let  $e \in E(T)$ , where  $T_1$  and  $T_2$  are the two components of T - e and  $T_i$  has order  $n_i$  (i = 1, 2). Assume, without loss of generality, that  $n_1 \leq n_2$ . Thus  $\operatorname{cn}(e) = n_1$ . Let s be a cyclic ordering of V(T), say  $s: v_1, v_2, \ldots, v_n, v_{n+1} = v_1$ . For each  $i \ (1 \leq i \leq n)$ , the edge e occurs at most once in the  $v_i - v_{i+1}$  path  $P_i$  of T. If e lies on  $P_i$ , then exactly one of  $v_i$  and  $v_{i+1}$  belongs to  $T_1$ . Since a vertex of  $T_1$  can occur as the initial or terminal vertex of a path  $P_i \ (1 \leq i \leq n)$  at most  $2 \operatorname{cn}(e)$  times, the desired result follows.

For the tree T of Figure 3,

 $\sum_{i=1}^{8} \operatorname{cn}(e_i) = 1 + 1 + 3 + 1 + 4 + 1 + 2 + 1 = 14.$ 

Thus by Lemma 4.4,  $h^+(T) \leq 28$ . However, for

 $s: v_1, v_9, v_2, v_8, v_3, v_7, v_5, v_6, v_4, v_1,$ 

we have d(s) = 28. Therefore,  $d(s) = 28 \le h^+(T)$  and so  $h^+(T) = 28$ . We now present a formula for  $h^+(P_n)$ .

**Proposition 4.5** For each  $n \ge 2$ ,

$$h^+(P_n) = \lfloor n^2/2 \rfloor.$$

**Proof.** Let  $P_n: v_1, v_2, \ldots, v_n$  and let

$$s: v_1, v_n, v_2, v_{n-1}, v_3, \ldots, v_{\lceil \frac{n+1}{2} \rceil}, v_1.$$

Then

$$d(s) = (n-1) + (n-2) + \dots + 1 + \left\lceil \frac{n-1}{2} \right\rceil \\ = \binom{n}{2} + \left\lceil \frac{n-1}{2} \right\rceil = \left\lceil \frac{n^2 - 1}{2} \right\rceil = \left\lfloor \frac{n^2}{2} \right\rfloor.$$

Hence  $h^+(P_n) \ge \left\lfloor \frac{n^2}{2} \right\rfloor$ .

To show that  $h^+(P_n) \leq \left\lfloor \frac{n^2}{2} \right\rfloor$ , we consider two cases, according to whether *n* is odd or *n* is even. Let  $e_i = v_i v_{i+1}$ ,  $1 \leq i \leq n-1$ .

Case 1. n is odd, say n = 2k + 1, where  $k \ge 1$ . Then

$$\operatorname{cn}(e_i) = \begin{cases} i & \text{if } 1 \leq i \leq k \\ n-i & \text{if } k+1 \leq i \leq 2k. \end{cases}$$

By Lemma 4.4,

$$h^{+}(P_{n}) \leq 2 \quad \sum_{i=1}^{n-1} \operatorname{cn}(e_{i}) = 2 \left[ \sum_{i=1}^{k} \operatorname{cn}(e_{i}) + \sum_{i=k+1}^{2k} \operatorname{cn}(e_{i}) \right]$$
$$= 4 \sum_{i=1}^{k} i = 4 \binom{k+1}{2} = 4 \binom{\frac{n+1}{2}}{2} = \frac{n^{2}-1}{2}.$$

Case 2. n is even, say n = 2k, where  $k \ge 1$ . Then

$$\operatorname{cn}(e_i) = \begin{cases} i & \text{if } 1 \leq i \leq k \\ n-i & \text{if } k+1 \leq i \leq 2k-1. \end{cases}$$

By Lemma 4.4,

$$h^{+}(P_{n}) \leq 2 \quad \sum_{i=1}^{n-1} \operatorname{cn}(e_{i}) = 2 \left[ \sum_{i=1}^{k} i + \sum_{i=k+1}^{2k-1} (n-i) \right]$$
$$= 2 \left[ \sum_{i=1}^{k} i + \sum_{i=1}^{k-1} i \right] = 2 \left[ 2 \sum_{i=1}^{k-1} i + k \right]$$
$$= 2 \left[ 2 \binom{k}{2} + k \right] = 4 \binom{\frac{n}{2}}{2} + n = \frac{n^{2}}{2}.$$

Thus, in each case,  $h^+(P_n) \leq \left\lfloor \frac{n^2}{2} \right\rfloor$ , producing the desired result.

If T is a tree of order n and T' is a tree obtained by adding a pendant edge to T, then  $cn_T(e) \leq cn_{T'}(e) \leq cn_T(e) + 1$  for every edge e of T. We now show that the upper bound is attained for at most half of the edges of T. With the aid of this fact, we will be able to establish a sharp upper bound for the upper Hamiltonian number of a graph in terms of its order.

**Lemma 4.6** Let T be a tree of order n, and let T' be a tree obtained by adding a pendant edge to T. Then there are at most (n-1)/2 edges e in T such that  $\operatorname{cn}_{T'}(e) = \operatorname{cn}_T(e) + 1$ .

**Proof.** For each  $e \in E(T)$ , let  $T_{1e}$  and  $T_{2e}$  be the two components of T-eand let  $n_{1e}$  and  $n_{2e}$  be the orders of  $T_{1e}$  and  $T_{2e}$ , respectively. Assume, without loss of generality, that  $n_{1e} \leq n_{2e}$ . Thus  $cn(e) = n_{1e}$ . Let  $e_0 = xy$ be an edge of T such that  $n_{2e_0} - n_{1e_0} \leq n_{2e} - n_{1e}$  for all edges e in T. Suppose that T' is obtained from T by adding the pendant edge uv at the vertex u of T. We show that the number of edges e in T such that  $cn_{T'}(e) = cn_T(e) + 1$  is at most (n-1)/2. Let  $T_1$  and  $T_2$  be the two components of  $T - e_0$  such that  $cn(e_0)$  is the order of  $T_1$  We may assume that  $x \in V(T_1)$  and  $y \in V(T_2)$ . For each  $e \in E(T)$ , let  $T'_{1e}$  and  $T'_{2e}$  be the two components of T' - e and let  $n'_{1e}$  and  $n'_{2e}$  be the orders of  $T'_{1e}$  and  $T'_{2e}$ , respectively. We may assume that  $n'_{1e} \leq n'_{2e}$ . We consider two cases.

Case 1.  $u \in V(T_2)$ . Let P be the y - u path in  $T_2$  (it is possible that y = u) as shown in Figure 4. Let  $e \in E(T) - E(P)$ . We consider two possibilities.

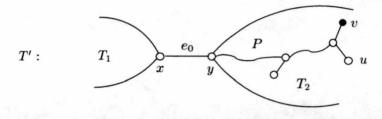


Figure 4: The tree T' in Case 1

Subcase 1.1.  $e \in E(T_1) \cup \{e_0\}$ . Then  $T'_{1e} = T_{1e}$ , while  $T'_{2e}$  is obtained by adding v and the edge uv to  $T_{2e}$ . Therefore,  $\operatorname{cn}_{T'}(e) = n'_{1e} = n_{1e} = \operatorname{cn}_{T}(e)$  for all  $e \in E(T_1) \cup \{e_0\}$ .

Subcase 1.2.  $e \in E(T_2) - E(P)$ . We show that  $\operatorname{cn}_{T'}(e) = \operatorname{cn}_T(e)$  in this subcase as well. Assume, to the contrary, that there exists  $f \in E(T_2) - E(P)$  such that  $\operatorname{cn}_{T'}(f) = \operatorname{cn}_T(f) + 1$ . Then  $T'_{1f}$  is obtained by adding the pendant edge uv to  $T_{1f}$ , while  $T'_{2f} = T_{2f}$ . Since x and v are connected in T' - f (by the path whose edge set is  $E(P) \cup \{e_0, uv\}$ ) and  $v \in E(T'_{1f})$ , it follows that  $T_1$  is a proper subgraph of  $T'_{1f}$ . Since  $T'_{1f}$  is obtained from  $T_{1f}$  by adding the pendant edge uv and  $uv \notin E(T_1)$ , it follows that  $T_1$  is a proper subgraph of  $T'_{2f}$  is a proper subgraph of  $T_2$ . This implies that  $n_{1e_0} < n_{1f} \leq n_{2f} \leq n_{2e_0}$  and so  $n_{2f} - n_{1f} < n_{2e_0} - n_{1e_0}$ , which is impossible.

Therefore, if  $e \in E(T)$  and  $\operatorname{cn}_{T'}(e) = \operatorname{cn}_T(e) + 1$ , then  $e \in E(P)$ . It remains to show that  $|E(P)| \leq (n-1)/2$ . Assume, to the contrary, that  $|E(P)| \geq n/2$ . Let

$$P': y = v_0, v_1, \ldots, v_{|E(P)|} = u, v$$

be the path obtained by extending P to v. Let  $f_0 = yv_1$  (see Figure 5). Then  $T_1$  and the path P' - y belong to different components in T' - y. Since the order of P' - y is  $|E(P)| + 1 \ge n/2 + 1$ , it follows that P' - y is a subgraph of  $T'_{2f_0}$ . Thus  $T_1$  is a proper subgraph  $T'_{1f_0} = T_{1f_0}$ . Since  $T_{2f_0}$  is a subgraph  $T_2$ , it follows that  $n_{1e_0} < n_{1f_0} \le n_{2f_0} \le n_{2e_0}$  and so  $n_{2f_0} - n_{1f_0} < n_{2e_0} - n_{1e_0}$ , which is impossible.

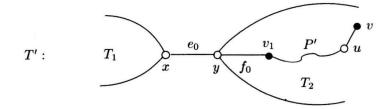


Figure 5: The path P' and the edge  $f_0$  in T' in Case 1

Case 2.  $u \in V(T_1)$ . Let Q be the u - x path in  $T_1$  (it is possible that u = x) as shown in Figure 6. Let  $e \in E(T) - (E(Q) \cup \{e_0\})$ . We now consider two subcases.

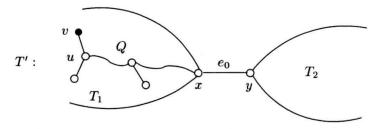


Figure 6: The tree T' in Case 2

Subcase 2.1.  $e \in E(T_1) - E(Q)$ . Then  $T'_{1e} = T_{1e}$ , while  $T'_{2e}$  is obtained by adding v and the edge uv to  $T_{2e}$ . Thus  $\operatorname{cn}_{T'}(e) = \operatorname{cn}_T(e)$  for all  $e \in E(T_1) - E(Q)$ .

Subcase 2.2.  $e \in E(T_2)$ . We show that  $\operatorname{cn}_{T'}(e) = \operatorname{cn}_T(e)$  in this subcase as well. Assume, to the contrary, that there exists  $f \in E(T_2)$  such that  $\operatorname{cn}_{T'}(f) = \operatorname{cn}_T(f) + 1$ . Then  $T'_{1f}$  is obtained by adding v and the edge uv to to  $T_{1f}$ , while  $T'_{2f} = T_{2f}$ . Since v and y is connected in T' - f (by the path whose edge set is  $E(Q) \cup \{uv, e_0\}$ ) and  $v \in V(T'_{1f})$ , it follows that  $T_1$  is a proper subgraph of  $T'_{1f}$ . Since  $T'_{1f}$  is obtained by adding vand the edge uv to to  $T_{1f}$  and  $uv \notin E(T_1)$ , it follows that  $T_1$  is a proper subgraph of  $T_{1f}$  and so  $T_{2f}$  is a proper subgraph of  $T_2$ . This implies that  $n_{1e_0} < n_{1f} \le n_{2f} < n_{2e_0}$  and so  $n_{2f} - n_{1f} < n_{2e_0} - n_{1e_0}$ , which is impossible.

Therefore, if  $e \in E(T)$  and  $\operatorname{cn}_{T'}(e) = \operatorname{cn}_T(e) + 1$ , then  $e \in E(Q) \cup \{e_0\}$ . We now consider the vertex  $e_0$ . If  $n_{1e_0} < n_{2e_0}$ , then  $T'_{1e_0}$  is obtained from  $T_1$  by adding the pendant edge uv and  $T'_{2e_0} = T_2$ , implying that  $\operatorname{cn}_{T'}(e_0) = \operatorname{cn}_T(e_0) + 1$ . If  $n_{1e_0} = n_{2e_0}$ , then  $T'_{2e_0}$  is obtained from  $T_1$ by adding the pendant edge uv and  $T'_{1e_0} = T_2$ , implying that  $\operatorname{cn}_{T'}(e_0) =$   $n_{2e_0} = n_{1e_0} = \operatorname{cn}_T(e_0)$ . Thus, there are two possibilities.

Case i.  $n_{1e_0} < n_{2e_0}$ . Therefore, if  $e \in E(T)$  and  $\operatorname{cn}_{T'}(e) = \operatorname{cn}_T(e) + 1$ , then  $e \in E(Q) \cup \{e_0\}$ . It remains to show that  $|E(Q) \cup \{e_0\}| \leq (n-1)/2$  or  $|E(Q)| \leq (n-3)/2$ . If  $|E(Q)| \geq (n-2)/2$ , then the order of Q is at least (n-2)/2 + 1 = n/2. On the other hand,  $n_{1e_0} < n_{2e_0}$  and so  $n_{1e_0} < n/2$ . However, Q is a subgraph of  $T_1$ , which is impossible.

Case ii.  $n_{1e_0} = n_{2e_0}$ . Therefore, if  $e \in E(T)$  and  $\operatorname{cn}_{T'}(e) = \operatorname{cn}_T(e) + 1$ , then  $e \in E(Q)$ . It remains to show that  $|E(Q)| \leq (n-1)/2$ . If  $|E(Q)| \geq n/2$ , then the order of Q is at least n/2 + 1. However, Q is a subgraph of  $T_1$  and the order of  $T_1$  is at most n/2, which is impossible.

An observation concerning trees that are not paths will also be useful.

**Lemma 4.7** If T is a tree of order  $n \ge 5$  that is not a path, then there exists an end-vertex v in T such that T - v is not a path.

For trees that are not paths, we can now establish an upper bound for the sum of the component numbers of its edges.

**Theorem 4.8** If T is a tree of order  $n \ge 4$  that is not a path, then

$$2\sum_{e\in E(T)}\operatorname{cn}(e)\leq \frac{n^2-4}{2}.$$

**Proof.** We proceed by induction on n. If n = 4, then  $T = K_{1,3}$  is the only tree that is not a path. Since  $h^+(K_{1,3}) = 6 = \frac{4^2}{2} - 2$ , the result holds for n = 4. Suppose that the result holds for all trees of order  $n - 1 \ge 4$  that are not paths. Let T be a tree of order  $n \ge 5$  that is not a path. By Lemma 4.7 there exists an end-vertex v in T such that T - v is not a path. Assume, without loss of generality, that  $E(T) = \{e_1, e_2, \ldots, e_{n-2}, e_{n-1}\}$  and  $E(T-v) = E(T) - \{e_{n-1}\}$ . Then  $\operatorname{cn}_T(e_{n-1}) = 1$  and by the induction hypothesis,

$$2\sum_{i=1}^{n-2} \operatorname{cn}_{T-v}(e_i) \le \frac{(n-1)^2 - 4}{2}.$$
(5)

It then follows by Lemma 4.6 that

$$2\sum_{i=1}^{n-1} \operatorname{cn}_{T}(e_{i}) = 2\sum_{i=1}^{n-2} \operatorname{cn}_{T}(e_{i}) + 2\operatorname{cn}_{T}(e_{n-1})$$
$$\leq 2\left[\sum_{i=1}^{n-2} \operatorname{cn}_{T-v}(e_{i}) + \frac{n-2}{2}\right] + 2\operatorname{cn}_{T}(e_{n-1})$$

If n is even, then by (5)

$$2\sum_{i=1}^{n-1} \operatorname{cn}_T(e_i) \leq \frac{(n-1)^2 - 4}{2} + (n-2) + 2$$
$$\leq \frac{(n-1)^2 - 5}{2} + n = \frac{n^2 - 4}{2}.$$

If n is odd, then by (5)

$$2\sum_{i=1}^{n-1} \operatorname{cn}_{T}(e_{i}) \leq 2\left[\sum_{i=1}^{n-2} \operatorname{cn}_{T-v}(e_{i}) + \frac{n-3}{2}\right] + 2$$
$$\leq \frac{(n-1)^{2}-4}{2} + (n-3) + 2 \leq \frac{n^{2}-5}{2}$$

Thus, in each case,  $2\sum_{i=1}^{n-1} \operatorname{cn}_T(e_i) \leq \frac{n^2-4}{2}$ , as desired.

As with the Hamiltonian number, if G if a connected graph of order  $n \ge 4$  and H is a connected spanning subgraph of G, then  $h^+(G) \le h^+(H)$ . Thus, the following result follows by Theorem 4.8.

**Corollary 4.9** Let G be a connected graph of order  $n \ge 4$  that is not a path. Then

$$h^+(G) \le \lfloor n^2/2 \rfloor - 2.$$

It then follows by Corollary 4.9 and Theorem 4.5 that there is no connected graph G of order  $n \ge 4$  having  $h^+(G) = \lfloor n^2/2 \rfloor - 1$ . The following is a consequence of Theorems 2.3, 3.2, and 4.8.

**Corollary 4.10** Let T be a tree of order  $n \ge 3$ . Then

$$2n-2 \le h^+(T) \le |n^2/2|$$
.

Moreover,

- (a)  $h^+(T) = 2n 2$  if and only if  $T = K_{1,n-1}$ ,
- (b)  $h^+(T) = |n^2/2|$  if and only if  $T = P_n$ .

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